

Algorithmic Model Theory — Assignment 3

Due: Friday, 6 May, 13:00

Exercise 1

We say that a domino system $\mathcal{D} = (D, H, V)$ admits a *horizontally periodic* tiling if we can find a tiling $\tau : \mathbb{N} \times \mathbb{N} \rightarrow D$ and an integer $h \geq 1$ such that for all $(x, y) \in \mathbb{N} \times \mathbb{N}$ we have $\tau(x, y) = \tau(x + h, y)$.

Prove or disprove that every domino systems \mathcal{D} which admits a horizontally periodic tiling also admits a periodic tiling.

Exercise 2

In this exercise we want to show that the model construction for FO^2 -formulae from the lecture is optimal in the following sense: in general it does not suffice to take, for each $\forall\exists$ -subformula, only two copies (instead of three) of the set P which consists of atomic 1-types which are realised at least twice in \mathfrak{A} .

Find a satisfiable FO^2 -sentence $\varphi = \forall x \forall y \alpha \wedge \forall x \exists y \beta$ where α, β are quantifier-free such that:

- no model of φ contains a king (i.e. $K = \emptyset$) and
- for every model \mathfrak{A} of φ there is no corresponding finite model over the universe $P \times \{0, 1\}$.

Exercise 3

In the lecture we saw that the class \mathcal{F} consisting of all FO-sentences of the form $\forall x \eta(x)$ (for quantifier-free η) which only contain unary function symbols is a conservative reduction class.

Prove that the more restricted class $\mathcal{F}_2 \subseteq \mathcal{F}$ consisting of sentences in \mathcal{F} that contain only two unary function symbols is also a conservative reduction class.

Hint: Transform sentences $\forall x \varphi$ with unary function symbols f_1, \dots, f_m into sentences $\forall x \tilde{\varphi} := \forall x \varphi[x/hx, f_i/hg^i]$ where h, g are fresh unary function symbols.

Exercise 4

Show that the following classes of FO-sentences, where R is a binary relational symbol and f is a unary function symbol, contain infinity axioms.

- $\exists x \forall y \forall z \eta(x, y, z)$, $\eta \in \text{FO}(\{f\})$ quantifier-free.
- $\forall x \exists y \forall z \eta(x, y, z)$, $\eta \in \text{FO}(\{R, f\})$ quantifier-free and without equality.
- $\forall x \exists y \eta(x, y)$, $\eta \in \text{FO}(\{f\})$ quantifier-free.
- The two variable fragment FO^2 extended by the counting quantifiers $\exists^{\leq n}$ for every $n \in \mathbb{N}$, where $\exists^{\leq n} x \varphi$ expresses that there are no more than n elements x that satisfy φ .