

## Algorithmic Model Theory — Assignment 4

Due: Monday, 14 November, 12:00

### Exercise 1

Let  $\tau_n$  be the vocabulary consisting of the unary relation symbols  $P_0, \dots, P_{n-1}, Q_0, \dots, Q_{n-1}$ .

Construct a formula  $\varphi_n \in \text{FO}(\tau_n)$  whose length is bounded by a polynomial in  $n$  such that every model of  $\varphi_n$  contains at least  $2^{2^n}$  elements.

*Hint:* Encode an  $(2^n \times 2^n)$ -grid in models of  $\varphi_n$ .

### Exercise 2

$\varepsilon$ -FO is the extension of FO by Hilbert's *choice operator* (also known as  $\varepsilon$ -operator). The syntax of  $\varepsilon$ -FO is given by the usual rules together with an additional  $\varepsilon$ -rule: If  $\psi$  is a formula, and  $x$  is a variable, then  $\varepsilon x\psi$  is a term (read "an  $x$  such that  $\psi$ ").

An interpretation for an  $\varepsilon$ -FO formula is given by an FO-interpretation  $(\mathfrak{A}, \mathcal{I})$  together with a choice function on the universe of  $\mathfrak{A}$ , i.e. a mapping  $F : \mathcal{P}(A) \rightarrow A$  such that  $F(X) \in X$  for all  $X \neq \emptyset$ . The value of a term  $\varepsilon x\psi$  is defined as  $F(\{a \in A : (\mathfrak{A}, \mathcal{I}) \models \psi(a)\})$ .

- Show that the quantifiers  $\exists$  and  $\forall$  can be expressed with the  $\varepsilon$ -operator.
- Construct an infinity axiom  $\varphi$  in  $\varepsilon\text{-FO}^2$  over the empty vocabulary, i.e.,  $\varphi$  contains only the  $\varepsilon$ -operator, two variables  $x$  and  $y$ , and equality, but neither relation nor function symbols.

### Exercise 3

Specify a conservative reduction from  $[\forall^2, (0, 1), (1)]$  to  $[\forall^2, (1), (0, 1)]$ , i.e., from the class of formulae with quantifier prefix  $\forall^2$  over a vocabulary containing only one binary relation and a unary function symbol to the class of formulae with the same quantifier prefix but over a vocabulary containing a unary relation and a binary function symbol. Note that equality may not be used in the formulae.

*Hint:* A binary relation  $P$  and a unary function  $f$  can be encoded using a unary relation  $Q$  and a binary function  $g$  in the following way:  $P = \{(a, b) : g(a, b) \in Q\}$ , and  $f(a) = g(a, a)$ .