Lehr- und Forschungsgebiet Mathematische Grundlagen der Informatik RWTH Aachen Prof. Dr. E. Grädel, F. Abu Zaid, W. Pakusa

## Algorithmic Model Theory — Assignment 12

Due: Monday, 23 January, 12:00

## Exercise 1

Show that the following classes of (undirected) graphs are definable in FPC.

- (a) Graphs having an odd number of edges. Hint: Handshaking lemma
- (b) Graphs having an even number of connected components.

## Exercise 2

Let  $\mathfrak{A}$  be a finite  $\tau$ -structure. We make the following convention: we interpret numerical tuples  $\bar{\nu} = (\nu_{k-1}, \ldots, \nu_1, \nu_0) \in \{0, \ldots, |A| - 1\}^k$  as numbers in |A|-adic representation, i.e. we associate the value  $\sum_{i=0}^{k-1} \nu_i |A|^i$  to each tuple  $\bar{\nu} \in \{0, \ldots, |A| - 1\}^k$ .

Show that the expressive power of FPC does not increase if we allow counting quantifiers of higher arity, i.e. formulas  $\#_{x_0x_1\cdots x_{k-1}}\varphi(x_0,\ldots,x_{k-1}) \leq (\nu_{k-1},\ldots,\nu_0)$  where in a structure  $\mathfrak{A}$  the value of  $\#_{x_0x_1\cdots x_{k-1}}\varphi(x_0,\ldots,x_{k-1})$  is the number of tuples  $\bar{a}$  such that  $\mathfrak{A} \models \varphi(\bar{a})$  (with respect to the encoding introduced above). For simplicity, you may only consider the case k = 2.

## Exercise 3

We denote by  $C_{\infty\omega}^k$  the k-variable infinitary logic with counting which is the extension of  $L_{\infty\omega}^k$  by all quantifiers of the form  $\exists^{\geq n} x$  with the intended semantics "there are at least *n* elements x such that". Similarly to the definition of  $L_{\infty\omega}^{\omega}$  we set  $C_{\infty\omega}^{\omega} := \bigcup_{k\geq 1} C_{\infty\omega}^k$ .

For illustration, consider the following sentence of  $C^1_{\infty\omega}$  defining the class of graphs with an even number of vertices

$$\bigvee_{n\in 2\omega} (\exists^{\geq n} x(x=x) \land \neg \exists^{\geq n+1} x(x=x)).$$

Show that FPC  $\leq C_{\infty\omega}^{\omega}$  (for formulas without free numerical variables). Hints:

• Construct for every formula  $\varphi(\bar{x},\bar{\nu}) \in \text{FPC}$  and  $n \in \omega, \bar{\nu} \in \{0,\ldots,n-1\}^k$  a formula  $\varphi^*_{n,\bar{\nu}}(\bar{x})$  which is equivalent to  $\varphi$  on structures of size n, i.e. for all  $\mathfrak{A}$  of size n we have

$$\mathfrak{A} \models \varphi(\bar{a}, \bar{\nu}) \text{ iff } \mathfrak{A} \models \varphi_{n, \bar{\nu}}^{\star}(\bar{a}), \text{ for all } \bar{a} \in A.$$

• For fixed point operators, adapt the construction from the proof showing  $FP \leq L^{\omega}_{\infty\omega}$ .

http://logic.rwth-aachen.de/Teaching/AMT-WS12/