Lehr- und Forschungsgebiet Mathematische Grundlagen der Informatik RWTH Aachen Prof. Dr. E. Grädel, F. Abu Zaid, W. Pakusa, F. Reinhardt

Algorithmic Model Theory — Assignment 2

Due: Monday, 4 November, 12:00

Note: - You may work on the exercises in groups of up to three students.

 Hand in your solutions at the end of the lecture or put them into the box at the institute.

Exercise 1

Let X be the set of relational FO-sentences of the form $\exists x_1 \dots \exists x_r \forall y_1 \dots \forall y_s \varphi$ where $r, s \in \mathbb{N}$ and φ is quantifier-free. Show that Sat(X) is decidable.

Hint: Show that each satisfiable sentence in X has a model with at most r elements.

Exercise 2

(a) Show that $[\forall^3 \exists, (0, \omega), (0)]_{=}$ is a conservative reduction class.

Hint: Use the same technique as in reduction from the domino problem to the $\forall \exists \forall \text{-class}$, but use a binary relation to describe the successor function.

(b) Show that this even holds in the absence of equality, i.e. show that $[\forall^3 \exists, (0, \omega), (0)]$ is a conservative reduction class.

Hint: Try to substitute equality by an appropriate congruence relation.

Exercise 3

Which of the following subclasses of CORNER-DOMINO are r.e. and which are co-r.e.? In each case prove your answer.

- (i) CORNER-DOMINO = { (\mathcal{D}, D_0) : there exists a tiling of $\mathbb{N} \times \mathbb{N}$ by \mathcal{D} with origin constraint D_0 }
- (ii) CDOMINO-PER = { (\mathcal{D}, D_0) : there exists a periodic tiling of $\mathbb{N} \times \mathbb{N}$ by \mathcal{D} with origin constraint D_0 }
- (iii) CDOMINO-NPER = { (\mathcal{D}, D_0) : there exists a non-periodic tiling of $\mathbb{N} \times \mathbb{N}$ by \mathcal{D} with origin constraint D_0 but no periodic one}
- (iv) CDOMINO-UNIQUE = { (\mathcal{D}, D_0) : there exists a unique tiling of $\mathbb{N} \times \mathbb{N}$ by \mathcal{D} with origin constraint D_0 }