Lehr- und Forschungsgebiet Mathematische Grundlagen der Informatik RWTH Aachen Prof. Dr. E. Grädel, F. Abu Zaid, W. Pakusa, F. Reinhardt

# Algorithmic Model Theory — Assignment 6

Due: Monday, 2 December, 12:00

## Exercise 1

To justify the definition of SO-HORN, show that the admission of arbitrary first-order prefixes would make the restriction to Horn clauses pointless. This extension of SO-HORN has the full power of second-order logic.

## Exercise 2

weak-SO-HORN is the set of all formulae of the form

$$QR_1 \dots QR_k \forall x_1, \dots \forall x_l \bigwedge_{1 \le i \le r} C_i.$$

The  $C_i$  are of the form  $B_1 \wedge \ldots \wedge B_n \to H$  where the  $B_i$  are either atoms or negated atoms with the restriction that the relations  $R_1, \ldots, R_k$  may only occur positively. That means weak-SO-HORN differs from SO-HORN in that fact that only atomic or negated atomic first order formulae are allowed in the clauses.

- (a) Show that on ordered structures weak-SO-HORN is strictly less expressive than SO-HORN. *Hint:* Show that for every weak-SO-HORN sentence  $\psi$  the class  $\{\mathfrak{A} : \mathfrak{A} \models \psi\}$  is closed under substructures.
- (b) Show that, however, on ordered structures with the additional successor relation and constants 0, *e* for the first and last element in the order weak-SO-HORN and SO-HORN are equally expressive.

*Hint:* show that on this domain weak-SO-HORN captures PTIME.

## Exercise 3

A finite game  $\mathcal{G} = (V, V_0, V_1, E)$  is called *k*-bounded if every  $v \in V$  has at most *k* outgoing edges. Give SO-HORN formulae  $\varphi_k(v)$  so that for all *k*-bounded games  $\mathcal{G}$  and all  $v \in V$ 

 $\mathcal{G} \models \varphi_k(v) \Leftrightarrow$  player 0 has no winning strategy from v

 ${\it Remark:}$  For a definition of finite games see chapter 2 of the mathematical logic 1 lecture notes.

## Exercise 4

An operator  $F : \mathcal{P}(A) \to \mathcal{P}(A)$  is called *inflationary* if  $F(X) \supseteq X$  for all  $X \subseteq A$ . Give examples for operators  $F : \mathcal{P}(A) \to \mathcal{P}(A)$  with the following properties:

- (i) F has a fixed point but no least one.
- (ii) F has a least fixed point but is not monotone.
- (iii) F is monotone but not inflationary.
- (iv) F is inflationary but not monotone.

http://logic.rwth-aachen.de/Teaching/AMT-WS13/