Lehr- und Forschungsgebiet Mathematische Grundlagen der Informatik RWTH Aachen Prof. Dr. E. Grädel, K. Dannert

Algorithmic Model Theory — Assignment 5

Due: Tuesday, 12 November, 10:30

Exercise 1

In the definition of SO-HORN, the clauses $\beta_1 \wedge \cdots \wedge \beta_m \to H$ are Horn-clauses only with respect to R_1, \ldots, R_m , since the β_i can be arbitrary FO(τ)-formulae (without R_i). We now consider a restriction SO-HORN^{*} of SO-HORN, where the β_i have to be atomic FO-fromulae. Prove or disprove that SO-HORN^{*} captures PTIME over ordered structures ($\mathfrak{A}, <$).

Exercise 2

(a) Describe the meaning of the LFP-sentence

$$\psi := \forall y \exists z F y z \land \forall y [lfp Ry. \forall x (Fxy \to Rx)](y)$$

in your own words and explain why it is an infinity axiom.

(b) Describe the meaning of the LFP-sentence

$$\forall x \forall y (Exy \to \neg [\operatorname{lfp} Pz. \ z = x \lor (\exists u \exists v ((Ezu \lor Euz) \land (Euv \lor Evu) \land Pv))](y)).$$

Exercise 3

- (a) Let E be a binary relation symbol and P a unary relation symbol. Give an LFP($\{E, P\}$)-formula expressing the following properties:
 - (i) There is an element from which no E-path contains a node in P.
 - (ii) There is an element from which there is an E-path containing infinitely many nodes in P.
- (b) Let GEN = (A, f, S, u), where $f : A \times A \to A$ is a binary function, S is a unary predicate and u is a constant. Give an LFP-fromula φ such that $\text{GEN} \vDash \varphi$ if and only if u is in the closure of S under f.

Exercise 4

Let $\varphi(x) = [\text{lfp } Ry. \exists x_1 \dots \exists x_m \psi(R, y, x_1, \dots, x_m)](x)$ with ψ quantifier-free. Prove that for all structures \mathfrak{A} (finite or infinite) and all $a \in A$, $\mathfrak{A} \models \varphi(a)$ if and only if $a \in R^{\omega}$, i.e. the fixed point is reached at (or before) the first limit ordinal.

http://logic.rwth-aachen.de/Teaching/AMT-WS19/

5 Points

8 Points

12 Points

5 Points